

# KAOS

For People Who Have Got Smart

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Last week we received a letter from Kerry Lourash, who had some how managed to see a copy of our newsletter in America, and he made some complementary remarks about it. Kerry sent us a very interesting program which is on page 6 of this month's newsletter. He would like to correspond with fellow ML hackers with a view to trading tips and information. Kerry's address is:

Kerry Lourash

U.S.A.

In view of the interest being shown in the use of computers in education, and as Tasmania is the leading state in Australia, in their use, we asked KAOS member, Graeme Reardon, who is the senior Maths/Science master at Cressy District High School, to write an article describing the TASNET system. His article is on page 13.

As David Wilson will not be able to attend the next meeting, there will not be a Forth meeting, but David Anear has been busily writing Forth programs and will be demonstrating his version of BREAKFORTH and his new R.G.B. driver.

The next meeting will be held on Sunday 29th August at 2pm at the Essendon Primary School, corner of Raleigh St and Nicholson St, Essendon. As this meeting is during the school holidays, we are not sure how many school children will be attending the morning session, but feel certain the very keen ones will be there, so would the early arrivers please bring their computers along as usual.

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## MORE MODIFICATIONS TO DISK BASIC

A few months ago, I showed how to create some free space within the BASIC interpreter. This month I will describe how I have used up some of this space. The routines described previously must have been implemented (unless you relocate the routines given here outside the areas normally used by DISK BASIC (0200 to 2300)).

One of my complaints regarding BASIC is the clumsiness of the facilities for interfacing with machine-code subroutines, especially in the area of passing parameters and returning results.

There are two basic methods:

POKE 8956, HI : POKE 8955, LO: Y=USR(X)

which allows a single integer parameter to be passed and a single integer result to be returned. The POKE addresses, of course, differ from ROM BASIC. It would make it too easy to transfer programs from ROM-BASIC to DISK-BASIC if they used the same addresses. The address of the routine must be known (in decimal) and the low and high parts calculated and POKEd.

The second method is specific to the disk system:

DISK!"GO 1234" which requires that the address be known in hexadecimal and does not permit any parameters to be passed or results to be returned.

In Disk BASIC, both the USR function and the DISK! operation cause pages 0 and 1 to be swapped (with 2F79 to 3178) producing an overhead of more than 21ms (at 1MHz) which removes some of the advantages of machine coded routines for high-speed applications. (The page swap can be by-passed in the case of the USR function. This strategy is not advised for the DISK! function.)

To overcome these problems, I 'invented' a CALL statement which replaces the NULL statement, a statement that is hardly ever used and can 'easily' be replaced. NULLX is equivalent to POKE21,X on a disk-system (POKE13,X on a ROM system).

There are two forms of the CALL statement implemented. The address of the machine code routine may be given by either a hexadecimal constant or a numeric expression (sorry, no strings allowed).

Parameters are 'easily' passed and results returned by CALLing suitable routines within BASIC. To illustrate this, the code below includes two, CALLable routines.

CALL\$17B6,P1,P2 (or CALL 6070,P1,P2) returns the address of P2 in P1. P1 should be a simple real variable. P2 may be an integer, string or real variable and may be simple or an array element.

CALL\$186C,X prints X on the current default device in hexadecimal. X should be between 0 and 255. The address of the CALLEd routine is assumed to be a normal BASIC expression unless it begins with a \$ in which case it must contain 4 hex digits.

The modifications to BASIC are:

0222	D0 17	(CALL-1)	replaces address of NULL routine
02C5	43 41	'CA'	(replaces 'NU')
17B6	20130E	JSR 0E13	check for comma
17B9	202E0F	JSR 0F2E	get address of first parameter
17BC	8596	STA 96	store address
17BE	8497	STY 97	
17C0	20130E	JSR 0E13	check for second comma
17C0	20130E	JSR 0F2E	get address of 2nd parameter
17C6	850A	STA 0A	swap A register and low in Y
17C9	A40A	LDY 0A	
17CB	201812	JSR 1218	convert integer to real
17CE	4CCB1A	JMP 1ACB	store result into (96,97) and exit
CALL			
17D1	C924	CMP #24	compare with '\$' hex address?
17D3	D010	BNE 17E5	no, decimal address instead
17D5	205519	JSR 1955	get two hex characters
17D8	851A	STA 1A	high byte of routine address
17DA	205519	JSR 1955	next two hex characters
17DD	8519	STA 19	low byte of address

17DF	20C000	JSR	00C0	get next character/token
17E2	6C1900	JMP	(0019)	jump to CALLED routine
17E5	4C6318	JMP	1863	get decimal address
1863	20B90C	JSR	0CB9	evaluate address of routine
1866	207216	JSR	1672	check range ( $0 \leq X \leq 65535$ ) and store in 19,1A
1869	6C1900	JMP	(0019)	jump to CALLED routine
186C	206C16	JSR	166C	evaluate expression and check range ( $0 \leq X \leq 255$ )
186F	8A	TXA		transfer value to A
1870	4C992D	JMP	2D92	print in hex and exit
1955	206419	JSR	1964	convert hex digit to binary
1958	0A	ASL	A	shift
1959	0A	ASL	A	to
195A	0A	ASL	A	top
195B	0A	ASL	A	nibble
195C	8519	STA	19	
195E	206419	JSR	1964	get 2nd hex digit
1961	0519	ORA	19	combine with first digit
1963	60	RTS		Return with result in A
1964	20C000	JSR	00C0	get next character
1967	2093FE	JSR	FE93	convert hex digit to binary using Monitor ROM
196A	10F7	BPL	1963	return if valid hex character
196C	4C1E0E	JMP	0E1E	invalid hex character, so report SN error

The following program demonstrates the use of the CALL statement with the two CALLable routines given above. The program takes a decimal value between 0 and 65535 and prints it out as a 4 character hexadecimal value.

```

5 PRINT"convert decimal to hex"
10 DISK!"ME 7000,7000"      set memory pointer to spare memory
20 A$="ZZ":X=0              define variables
30 CALL6070,X,A$            get location of string descriptor
40 POKEX+1,0:POKEX+2,112    point A$ to $7000
45 INPUT"Decimal";Z
47 Z1=INT(Z/256)            get top byte
48 Z=Z-Z1*256              get low byte
50 DISK!"IO,10              set default output to memory
60 CALL$186C,Z1             output top byte in hex (to memory)
70 CALL$186C,Z              output low byte in hex
80 DISK!"IO,02              set default output back to screen
90 POKEX,4                  set LEN(A$) to 4
100 PRINT A$                print result
110 GOTO 10                 reset pointers and get next value

```

This program is unnecessarily complicated in order to demonstrate the use of both CALLable routines. The program really only requires lines 45,47,48,60 and 70, which must make it the shortest BASIC program for the Challenger to do decimal to hex conversion.

Rodney Eisfelder

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#### FOR SALE

Programmable Sound Generator Board with construction details, cost \$15.00, or available as a kit or fully assembled and tested. Contact Ron Kerry at the KAOS meeting or phone or write to

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Siemens Model 100, Friction Feed, Auto switch off mechanism, Tape Reperforator and Tape transmitter, attachments, governed to 75 BWDS. V.G.C. Price \$285.00  
W. Babb

# Superboard

NEWSLETTER OF THE OHIO SUPERBOARD USER GROUP, 146 YORK ST., NUNDAH, QLD. 4012.

Some letters I received about the Basic 4 Eprom offered last month showed that readers were unclear about just which features of the original Basic 4 were deleted to make way for the new routines. NO useful features have been deleted. The routines replace some obsolete code left over from the days when the Basic was in RAM, as well as some unnecessary text.

This program, to suit owners of Cegmon monitor, will enable the screen and window clear keys CTRL Z and CTRL SHIFT N, direct from the keyboard. The cursor control keys, CTRL J,K,L, and M are also enabled. Normally, these are ?CHR\$(x) functions.

```
10 DATA 96,32,70,251,44,3,2,16,1,96,201,0,208,7,162,18,160,0
20 DATA 76,128,162,201,13,240,2,201,7,240,227,201,32,144,4
30 DATA 201,125,144,219,76,155,255
40 FOR R=6150 TO 6189:READ Z:POKE R,Z:NEXT:POKE 536,7:POKE 537,24
```

-----  
From Bob Best comes this neat mod to fix that Break Key problem. First, buy an 18¢ rubber grommet Dia. 17MM, thickness 5MM, hole 8MM from your hardware shop. Gently pry up the Keycap, place the grommet over the barrel, and restore the cap. Now a brush or accidental keypress will have no effect. You will need a firm push to operate the reset function.  
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## Hardware Review - EPROM EXTENDER

Eprom Extender is a bare 48mm X 120mm single sided printed circuit board, etched and drilled to accomodate three Eproms and a LS138 decoder I.C.

You solder in a wirewrap socket, and the pins plug into the Monitor Rom socket in the Superboard, which must be configured for a 2716 Eprom. The Monitor Eprom then plugs into the wirewrap socket. Three other connections are made from the Eprom Extender board to the Superboard, for decoding purposes, and the board is then strapped to provide a range of addresses for the two 2716 or 2732 Eproms which fit in the remaining two places on the Extender.

The review board was very well made, with no inter-track shorts or breaks, and was coated with protective resin. The instructions explained the wiring needed reasonably well, though a diagram showing the location of the chips on the Superboard would assist the owners without layouts and circuit diagrams. On the Eprom Extender, pin 18 is tied to Zero volts, so the chips run in the power-up mode and will operate at 2MHz.

To sum up:- A neat and simple way of obtaining an extra 4 or 8k of ROM.  
The Eprom Extender is available from Bert Patterson,  
Cost is \$8.50 plus \$1 for Pack & Post.

-----  
From Peter Wiseman comes this simple routine, when a few quick decimal to hex conversions are needed.

```
10 DATA 32,8,180,132,90,133,89,76,0,254
20 FOR R=565 TO 574:READ Z:POKE R,Z:NEXT:POKE 11,53:POKE 12,2
```

Run the program, then put the number you want to be converted to hex in the argument of the USR function. eg X = USR(65025) The program will do the job, then jump to monitor mode. Address \$0059 contains the high byte (FE)  
Address \$005A " " low " (01)

# Superboard

## Software Review - COMPACT

Compact is a machine code utility occupying \$1E00 - \$1FFF or \$3E00 - \$3FFF. Both versions are supplied on the tape.

As the name suggests, Compact is used to pack up Basic programs into the minimum number of lines, deleting lines that have been compacted, and no longer required. It does the job well, and quickly, a 5k Basic program compacting in under a minute. The only thing that I could get it to do wrong was to put two colons together on a line, which caused no problems in the subsequent running of the program.

The instructions are short, but adequate, giving directions on how to solve problems associated with Basic 5 statements in programs. The routine actually needs from \$1C00 to run, and you have to manually protect it from being overwritten by Basic before you load - something easy to forget. The program could have been made self-protecting by adding A9 00 85 85 A9 1C 85 86 4C 74 A2 before \$1E00, and making the routine self start from \$1DF5.

To sum up:- An excellent and useful utility.

Compact is available from Premier Publications, 208 Croydon Rd., Anerley, London SE20 7YX, priced at 8 Pounds + postage. The OSUG Library has one for Library members who would like to assess it personally. 27¢ + 40¢ stamps + label please.

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## Firmware Review - Character Sets

The product is an Eprom which can replace the OSI character generator, or be piggybacked onto it. This can give alternative character sets, either switch selectable or software selectable. The Eprom supplied can be a 2716 or a 2732. In both cases, modification of the computer board is very easy.

The following options are available:-

- (1) Standard Set. Available as one half of a 2732 Eprom only. For those who don't like to piggyback, and want an original set.
- (2) Enhanced Set. Basically similar to OSI, but some of the rarely used characters have been replaced. All letters and numbers have been shortened by one pixel, making them infinitely easier to read on the screen. Small letters have descenders, which also improves the presentation. Added characters include black and matching white chess pieces, helicopters, cars and aircraft in all directions, 3/4 blocks for double resolution graphing, a parachute, and random characters and others to produce explosions and smoke or dust effects. Available as a 2716 or part of a 2732.
- (3) Hi-Res Set. As in Practical Electronics, April 1982, pages 41-44. These give 4x resolution and are especially suitable for 24 x 48 or 32 x 64 screens. Available as 2716 or part of a 2732.

By piggybacking a 2732 onto your OSI character generator, you can have three character sets available at the flick of a switch. Yet to come is a scientific set with gaming characters replaced by various symbols and subscripts etc. New graphics charts and full documentation are provided with each chip. 2716 costs \$12.80 and 2732 costs \$17.80 including packing and post. Send your details to Bernie Wills,

Ed Richardson.

OSI GOFAST  
*by Kerry Lourash*

GOFAST is a one-page machine language program designed to speed up GOTO and GOSUB commands. It implements labeled GOTOs and GOSUBs that can be renumbered and edited. With GOFAST, BASIC lines are found from 25% to over 300% faster. If you have a BASIC-in-ROM ClP or an OSI computer with a CTRL C vector in RAM, you can use this program.

A Little Background:

When a GOTO command is executed in BASIC, this is what happens: The number of the target line is converted to a two-byte hex number. The high byte of this number is compared to the high byte of the current line number (the line that contains the GOTO). If the high byte of the target line is greater than the high byte of the current line number, BASIC starts searching for the line at the next line after the GOTO. Otherwise, the search starts at the beginning of the program. The search is done by comparing the target line number with the line number stored at the start of every line. If a comparison fails, BASIC finds the location of the next line by using the next-line pointer and does another comparison.

A LINE OF BASIC:

/	00	/	53	/	03	/	0A	/	00	/	CODE
NULL			(\$0353)				(10)				
			NEXT LINE				LINE #				
			POINTER								

When the target line is found, its address is loaded into the parser pointer (\$C3,C4) and BASIC resumes program execution at the target line.

GOSUBs are handled a little differently. Firstly, the stack is checked to see if there is enough room to accomodate the information that must be stored. Next, the contents of the parser pointer (\$C3, C4), the current line number (\$87, 88), and a GOSUB token (\$8C) are pushed onto the stack. Counting the return address to the BASIC execution loop (which was already on the stack), seven bytes are stored. Finally, BASIC calls the GOTO routine, which finds the target line and resumes program execution.

A Faster Way:

BASIC searches for the target line every time a GOTO or GOSUB is encountered. This method is fine for occasional use. When the search is repeated thousands of times in a loop, however, the program slows down dramatically. Why does BASIC use this method? My guess is that this way is most memory-efficient.

Suppose BASIC used two-letter labels for target lines, with the address of each label in an indexed table. The table would occupy 26X26X2 bytes of RAM! If the labels were stored in the order of their occurrence in the program, memory usage might be less than that of the indexed table. Retrieval would be slower, however, because a comparison search would be necessary to find the address!

So, what is left? POKEing the hex address of the target line number after the GOTO or GOSUB token would be both time and memory efficient. Unfortunately, the OSI ASCII SAVE and LOAD routines couldn't handle those bytes. Converting the hex address to ASCII and storing it after the token would be better, but both of these methods would make the program logic undecipherable for us poor humans.

The GOFAST Way:

GOFAST uses one-letter labels (A-Z). This strategy keeps the label table small. More letters may be used in a label, but only the first letter is read. For a GOTO, the command is: #TO A. I use #>A for a GOSUB, but the ">" can be replaced with a different label if you wish. The target line is labeled with: REM#A. This label is placed in the line preceding the target line. You may use a separate line for the label or combine it with another line.

Since there are only 26 possible labels, it is advisable to use labels only for GOTOs or GOSUBs that are called frequently.

While running a program, CTRL C is disabled. The CTRL C vector at \$021C, 021D is changed to point to GOFAST. GOFAST contains a streamlined version of the BASIC execution loop that resides at \$A502-A5FE.

A short BASIC program is used to load GOFAST. It lowers the top-of-RAM pointer (MEMTOP) one page to protect GOFAST and sets the USR vector to the initialization routine.

Here's a short program to test the speed of the GOFAST routine:

```
30 X=USR (X):REM CALL INIZ. ROUTINE
40 FOR I=1 TO 10000
50 #TO LOOP: REM#LOOP
60 NEXT
```

Note that the target REM causes a jump to line 60, not line 50.

To convert the program to stock BASIC, change:

```
30 PRINT
50 GOTO 60
```

Put ten or twenty BASIC lines (1 PRINT, 2 PRINT, etc.) in line numbers 1-29. The stock program's execution time will increase 200-300%, but the GOFAST program will not slow down.

#### BASIC V.S. GOFAST:

Here is a list of the BASIC execution loop's functions:

- 1 CTRL C check
- 2 Copy parser pointer into CONT pointer (\$8B, 8C).
- 3 Check for null (start of line).
- 4 Check for a colon (start of statement).
- 5 If null or colon not found, print "SN ERROR".
- 6 Check next-line pointer to see if program ends.
- 7 If no more lines, to immediate mode.
- 8 Put current line number into current line pointer (\$87, 88).
- 9 Update parser pointer to start of line.
- 10 Call parser (\$00BC) to get first character of line.
- 11 Call BASIC execution routine (\$A5FF).
- 12 Loop to #1.

GOFAST's version of the loop eliminates #1, 2, and 5. The net result is that GOFAST can check for a "#" at the start of every statement and still run faster than stock BASIC! I wonder how much faster BASIC would run with most of the error checking removed? For those speed freaks among us, steps 6, 7, and half step 8 could be removed if the END command is used and there are no bugs in the program to be run.

See GOFAST Run:

GOFAST consists of three parts: First, a 52 byte table that holds the addresses of the 26 labels, located just above MEMTOP. Second, an initialization routine that is run at the start of the user's program. It searches the user's program for labels and stores their locations in the table. Third, an operation routine that executes the #TO # > commands.

Since the initialization routine fills the table with zeroes before it searches for labels, any call to a non-existent label will send the program to warmstart and the immediate mode.

#### Side Effects:

Every program has its weak points. GOFAST doesn't do a foolproof job of error checking when it loads addresses into the label table, or when it looks them up. As a result, addresses may be retrieved or stored in the GOFAST program area, if the correct command format is not used. The effects will be unpredictable and might crash the program. I highly recommend the standard practice of SAVEing a program at intermediate stages of development to be safe.

If a BASIC program is terminated with an END command, GOFAST will not reset the CTRL C vector. After running such a program, you will find that the LIST command doesn't work the first time it is used.

*Next page please*

GOFAST's END routine resets the CTRL C vector. C2P owners with a RAM CTRL C vector should change LDA #9B to LDA #F1. In the BASIC load program, change 155 to 241 in line 30250.

Many thanks to Earl "The Pearl" Morris for bug-catching and good advice.

A

```

1 0000 ;
10 0000 ; BASIC GOFAST MOD.
20 0000 ; BY KERRY LOURASH
30 0000 ;
40 0000 ; FORMAT: > A (GOSUB)
50 0000 ;      > TO B (GOTO)
60 0000 ; TARGET: REM#A
70 0000 ;      REM#B
80 0000 ;
90 0000 ; CTRLC=$021C    CTRL C VECTOR
100 0000 ; INDEX=$5E      TEMP. STORAGE FOR TABLE INDEX
110 0000 ; MEMTOP=$85      TOP OF RAM (+1)
120 0000 ; PARSER=$00BC  PARSER SUBROUTINE
130 0000 ; POINT=$AA       POINTER FOR INIT. ROUTINE
140 0000 ; START=$79      START OF BASIC WORKSPACE
150 0000 ;
160 0034 ; $=52        MAKE ROOM FOR TABLE
170 0034 ;
180 0034 ; INITIALIZATION ROUTINE
190 0034 ;
200 0034 A992 VECTOR LDA #ENTER*256/256 CHANGE CTRL C VECTOR
210 0036 8D1C02 STA CTRLC
220 0039 A900 LDA #ENTER/256
230 003B 8D1D02 STA CTRLC+1
240 003E ;
250 003E A033 ZROTLB LDY #51    ZERO LABEL TABLE
260 0040 A900 LDA #0
270 0042 9185 Z1 STA (MEMTOP),Y
280 0044 88 DEY
290 0045 30FB BPL Z1
300 0047 ;
310 0047 A579 SEARCH LDA START    SEARCH FOR REM LABELS
320 0049 85AA STA POINT
330 004B A57A LDA START+1
340 004D 85AB STA POINT+1
350 004F A003 NEXLIN LDY #3
360 0051 C8 NEXCHR INY
370 0052 B1AA LDA (POINT),Y
380 0054 F02D BEQ FIXLIN
390 0056 C98E CMP #8E    IS CHAR A REM?
400 0058 D0F7 BNE NEXCHR
410 005A C8 N1 INY
420 005B B1AA LDA (POINT),Y
430 005D C923 CMP #8    IS CHAR A "*" ?
440 005F 90F9 BCC N1    IS CHAR A SPACE?
450 0061 D020 BNE FIXLIN    IF # NOT FOUND, TO NEXT LINE
460 0063 C8 N2 INY
470 0064 B1AA LDA (POINT),Y    GET LABEL
480 0066 38 SEC
490 0067 E941 SBC #41
500 0069 90FB BCC N2    TRY AGAIN IF CHAR < #41
510 006B 0A ASL A    DOUBLE THE INDEX
520 006C 855E STA INDEX    SAVE THE INDEX
530 006E A000 STORE LDY #0    GET ADDRESS OF NEXT LINE
540 0070 B1AA LDA (POINT),Y
550 0072 E900 SBC #0    SUBTRACT ONE FROM LO BYTE
560 0074 AA TAX
570 0075 C8 INY    GET HI BYTE OF ADDRESS
580 0076 B1AA LDA (POINT),Y
590 0078 E900 SBC #0    DECREMENT IF NECESSARY
600 007A A45E LDY INDEX    STORE ADDRESS IN TABLE
610 007C C8 INY
620 007D 9185 STA (MEMTOP),Y
630 007F 88 DEY
640 0080 3A TXA
650 0081 9185 STA (MEMTOP),Y
660 0083 ;
670 0083 A000 FIXLIN LDY #0    SET POINT TO START OF LINE
680 0085 B1AA LDA (POINT),Y    GET LO BYTE OF ADDRESS
690 0087 AA TAX
700 0088 C8 INY
710 0089 B1AA LDA (POINT),Y    GET HI BYTE OF ADDRESS
720 008B 86AA STX POINT
730 008D 85AB STA POINT+1
740 008F D0BE BNE NEXLIN    IF HI BYTE <> 0, BRANCH
750 0091 60 RTS    RETURN TO BASIC
760 0092 ;
770 0092 ; OPERATION ROUTINE
780 0092 ;
790 0092 68 ENTER PLA    PULL SUB CALL FROM STACK
800 0093 68 PLA
810 0094 A000 LDY #0    FIRST CHAR OF LINE A NULL?
820 0096 B1C3 LOOP LDA ($C3),Y
830 0098 D01A BNE COLON    NO, MUST BE A COLON
840 009A AC02 NULL LDY #2    CHECK HI BYTE OF NEXT LINE
850 009C B1C3 LDA ($C3),Y    ADDRESS
860 009E F051 BEQ END    END OF BASIC PROGRAM

870 00A0 C8 INY    SAVE CURRENT LINE NUMBER
880 00A1 B1C3 LDA ($C3),Y
890 00A3 9587 STA $87
900 00A5 C8 INY
910 00A6 B1C3 LDA ($C3),Y
920 00A8 9588 STA $88
930 00AA 93 TYA    MOVE PARSER POINTER TO START
940 00AB 18 CLC    OF CODE
950 00AC 65C3 ADC $C3
960 00AE 65C3 STA $C3
970 00B0 9002 BCC COLON
980 00B2 E6C4 INC $C4
990 00B4 20BC00 COLON JSR PARSER    IS FIRST CHAR A "*" ?
1000 00B7 C923 CMP #8
1010 00B9 F007 BEQ GOFAST    YES, TO GOFAST
1020 00BB 20FFA5 JSR $A5FF    DO COMMAND
1030 00BE A000 LDY #0
1040 00C0 F0D4 BEQ LOOP    GO TO START OF EXEC. LOOP
1050 00C2 ;
1060 00C2 20BC00 GOFAST JSR PARSER    GET TOKEN, IGNORE SPACES
1070 00C5 AA TAX
1080 00C6 20BC00 JSR PARSER    GET LABEL, IGNORE SPACES
1090 00C9 E941 SBC #41    CALCULATE INBEX
1100 00CB 0A ASL A
1110 00CC AB TAY
1120 00CD E09D CPX #9D    IS TOKEN= "TO"?
1130 00CF F015 BEQ GOTO
1140 00D1 ;
1150 00D1 A9A5 GOSUB LDA #A5    PUSH RETURN ADDRESS
1160 00D3 48 PHA    ($A5C1)
1170 00D4 A9C1 LDA #C1
1180 00D6 48 PHA
1190 00D7 A5C4 LDA $C4    PUSH PARSER POINTER
1200 00D9 48 PHA
1210 00DA A5C3 LDA $C3
1220 00DC 48 PHA
1230 00DD A588 LDA $88    PUSH CURRENT LINE NUMBER
1240 00DF 48 PHA
1250 00E0 A587 LDA $87
1260 00E2 48 PHA
1270 00E3 A98C LDA #8C    PUSH GOSUB TOKEN
1280 00E5 48 PHA
1290 00E6 ;
1300 00E6 B185 GOTO LDA (MEMTOP),Y    GET ADDRESS FROM TABLE
1310 00E8 85C3 STA $C3    PUT IT IN PARSER POINTER
1320 00EA C8 INY
1330 00EB B185 LDA (MEMTOP),Y
1340 00ED 85C4 STA $C4
1350 00EF D0A9 BNE NULL    HI BYTE OF ADDRESS= ZERO?
1360 00F1 ;
1370 00F1 A99B END LDA #9B    RESTORE CTRL C VECTOR
1380 00F3 8D1C02 STA CTRLC
1390 00F6 A2FF LDX #FF
1400 00F8 8E1D02 STX CTRLC+1
1410 00FB CA DEX
1420 00FC 9A TXS
1430 00FD 4C74A2 JMP $A274    TO WARMSTART

LIST
30000 REM MOVE MEMTOP DOWN ONE PAGE
30010 A=PEEK(133):B=(PEEK(134)-1):POKE134,B
30020 REM CALCULATE START ADDRESS
30030 C=A+52+B*256
30040 REM POKE GOFAST INTO MEMORY
30050 FORI=0TO203:READ N:POKE C+I,N:NEXT
30060 REM CALC. ADDRESS OF ENTER ROUTINE
30070 REM AND POKE IT INTO INIT. ROUTINE
30080 E=A+146:IF E>256 THEN E=E-256:H=1
30090 H=H+B:POKE C+1,E:POKEC+6,H
30100 REM SET USR VECTOR TO INIT. ROUTINE
30110 E=A+52:IF E>256 THEN E=E-256:B=B+1
30120 POKE 11,E:POKE 12,B
30130 DATA169,128,141,28,2,169,2,141,29,2,160,51,169,0,145
30140 DATA133,136,16,251,165,121,133,170,165,122,133,171,160,3,200
30150 DATA177,170,240,45,201,142,208,247,200,177,170,201,35,144,249
30160 DATA208,32,200,177,170,56,233,65,144,248,10,133,94,160,0
30170 DATA177,170,233,0,170,200,177,170,233,0,164,94,200,145,133
30180 DATA136,138,145,133,160,0,177,170,170,200,177,170,134,170,133
30190 DATA171,208,190,96,104,104,160,0,177,195,208,26,160,2,177
30200 DATA195,240,81,200,177,195,133,135,200,177,195,133,136,152,24
30210 DATA101,195,133,195,144,2,230,196,32,188,0,201,35,240,7
30220 DATA32,255,165,160,0,240,212,32,188,0,170,32,188,0,233
30230 DATA65,10,168,224,157,240,21,169,165,72,169,193,72,165,196
30240 DATA72,165,195,72,165,136,72,165,135,72,169,140,72,177,133
30250 DATA133,195,200,177,133,133,196,208,169,169,155,141,28,2,162
30260 DATA255,142,29,2,202,154,76,116,162

```



## ANOTHER R.G.B. DRIVER BOARD

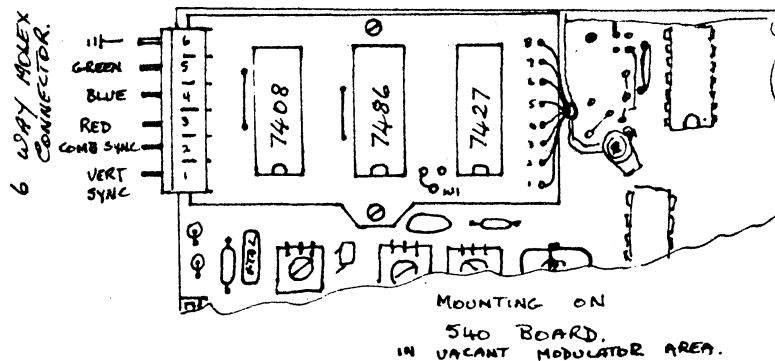
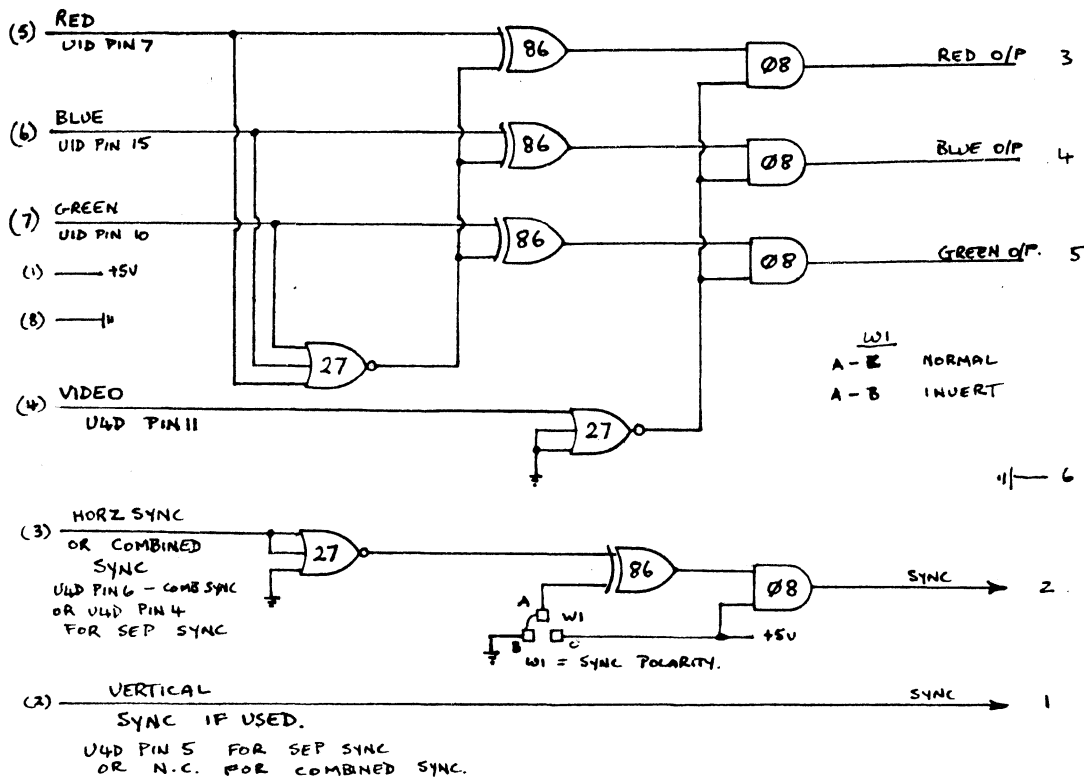
This little circuit has the following advantages over other designs:

1. OHIO compatible colours
2. Simple circuitry
3. Can drive modified T.V. or RGB monitor (see Vol.2 No.2 page 2)
4. Combined sync or separate sync O/P
5. By changing O/P IC open collector or inverse O/P can be selected
6. W1 used for inverting sync with respect to O/P sync's
7. Can be interfaced to Superboard series II, 540 Board in C4, C8P systems and TASAN Video board.

The board layout shown is for the C4/C8 system.

The RGB Board plugs into the J5 connector on the TASAN Video board with similar connections for the Superboard.  
If you are interested in the circuit boards, let me know and I will get some made.

David Anear



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## HOW TO COOK CHIPS BY BYTA BITROSE

Making unEproms into Eproms. Get out the frypan and set it on 175 C. (If cooking American chips, set it on 350 F). Put your unerased chips in the pan and leave them for about 20 minutes. Do not baste them as this will make them greasy chips. Remove them and place them on a plate to cool before letting your eraser eat the bytes.

This month, before looking at the instructions available on the 6502, we will cover the stack and the status register in a little more detail.

The STACK is an area of memory organised on a LAST IN, FIRST OUT basis. The operation of placing information on the stack is termed a PUSH. A PULL operation retrieves the information. We can simulate the operation of the stack in BASIC by treating it as a dimensioned array.

```
DIM STack (255):SP=255:REM SP=stack pointer:A=accumulator
PUSH:ST(SP)=A:SP=SP-1
PULL:SP=SP+1:A=ST(SP)
```

If you follow this through you will see information is being 'stacked' just like a pile of plates. As information is placed on the stack, the stack pointer is automatically updated to point to the next position on the stack. Likewise, before a PULL occurs, the stack pointer is updated to point to the last entry placed on the stack. In this way, items added to the stack can be accessed sequentially, but in reverse order. The architecture of the 6502 is such that the stack is located at \$0100-01FF (page 1).

The STATUS REGISTER comprises seven one bit 'flags' and an eighth bit which is not implemented. The flags are set to 1 for a true condition or cleared to zero for false. The flags are provided in order to:

- (a) make decisions based on the contents of a register or memory location.
- (b) indicate the current condition of the microprocessor.

The flags divide into two groups; four are known as status flags and the other three as condition flags. The STATUS flags are called carry (C), negative (N), zero (Z) and overflow (V).

In many operations, the CARRY can be considered to be a ninth bit on the accumulator. Its function is to signal when a carry or borrow is required as the result of an arithmetic operation.

The NEGATIVE (or sign) bit is set or cleared according to the state of the most significant bit after an operation on a register or memory. The CPU operates on the convention that a number is flagged as negative if the most significant bit is set, ie. signed binary arithmetic is assumed.

The ZERO flag is set to true (1) if the result of an arithmetic or boolean operation yielded an all zero result.

The OVERFLOW flag is seldom referenced in programming. It indicates whether the result of a twos' complement arithmetic operation is too big to be represented in one byte.

The CONDITION flags are called interrupt disable (I), break (B) and decimal mode (D).

The 6502 can operate in both binary and binary coded decimal (BCD) modes. In BCD mode, the maximum value that can be expressed in one byte is 99 (9x10+9 each represented in 4 bits). The carry is set if an operation causes a negative to exceed this value. In this way, the CPU simulates decimal arithmetic. The DECIMAL flag indicates which mode is currently in operation.

Most microprocessors can be interrupted while they are in mid program and forced to jump to another program. This facility is used by peripherals to let the CPU know they need attention rather than have the CPU constantly checking whether a peripheral device is ready. The 6502 has three levels of interrupt:

- non maskable
- maskable
- software

The BREAK key is an example of the way a non maskable interrupt (NMI) works; ie. it is always obeyed. The INTERRUPT DISABLE is the means used by the CPU to determine whether it is allowed to process a request by a peripheral for service. If interrupt disable is set the request will be ignored. Note that this flag is set by the CPU whenever any of the 3 types of interrupt occurs.

*Next page please*

The BREAK flag is set when a software interrupt (SWI or BREAK, the instruction not the key) is performed. Since NMI and SWI are treated the same way by the 6502, the programmer will decide what type of interrupt is being processed by examining the break flag.

The positions of the flags within the status register are:

Bit 0 C carry flag	Bit 4 B break flag
Bit 1 Z zero flag	Bit 5 unused (always 1)
Bit 2 I interrupt disable	Bit 6 V overflow flag
Bit 3 D decimal mode	Bit 7 N negative flag

The 6502 instructions can be divided into four distinct categories:

1. DATA TRANSFER instructions move data between two registers or between a register and memory. Some transfers involve arithmetic or logic operations. This group includes instructions such as LOAD, STORE, TRANSFER, ADD, SUBTRACT, AND, OR, EXCLUSIVE OR (EOR or XOR) and COMPARE. PUSH and PULL are also included in this category.
2. OPERATIONAL instructions differ from the former group in that while the contents of a register changes as a result of the operation, no data transfer takes place. Operational instructions include INCREMENT, DECREMENT, CLEAR, SET, BIT, TEST, SHIFT, and ROTATE.
3. CONTROL and BRANCH instructions are of two types: conditional and unconditional. Conditional branches test the state of one of the status registers and branch if the specified condition is met. The conditional branch works like the IF...THEN GOTO construction in BASIC. There are eight BRANCH instructions corresponding to each of the two states of the status flags. The unconditional branches include JUMP and BREAK.
4. SUBROUTINE LINKING gives the facility to call procedures located in another part of memory and to return to the same place where processing left off. These include JUMP to subroutine and RETURN from subroutine.

The 6502 offers 56 different instructions. To increase programming flexibility, there are 13 different methods of accessing information. Some instructions are available in as many as eight of these addressing modes.

In later articles, we will start to explore the various instructions and addressing modes.

David Dodds

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DEAR PAUL,

Q. What mods are necessary to change the CPU speed from 1 MHz to 2 MHz?

A. I suggest that you look back to Vol.1 No.8 of the KAOS newsletter and on page 9 you will find full circuits for accomplishing this remarkable feat.

Q. What is the difference between standard monitor and extended monitor?

A. The standard, 65V monitor supplied with the Superboard etc. has very minimal machine code programming facilities. The extended monitor is designed to allow complex machine code operations, it allows you to display a block of memory, edit memory, disassemble memory (display memory by showing you the 6502 opcodes) set breakpoints (points in your machine code program that cause control to be transferred to the extended monitor). The extended monitor is available on cassette or in EPROM, and comes standard with OS65D disks.

Q. What are the specific functions of FRE(X) and NULLX?

A. FRE(X) is a function which returns the amount of free memory, in bytes, it is usually used in the form PRINT FRE(X) or Y=FRE(X). NULLX is a statement which causes X null characters (CHR\$(0)) to be output after each carriage return, when data is being sent to the ACIA, ie. When you are saving a program to cassette. It is used when you have a device which takes a while to process a carriage return, eg. a teletype or a printer on the serial port.

Q. What is the @ command and how do you use it?

A. The @ command in BASIC simply deletes the current line.

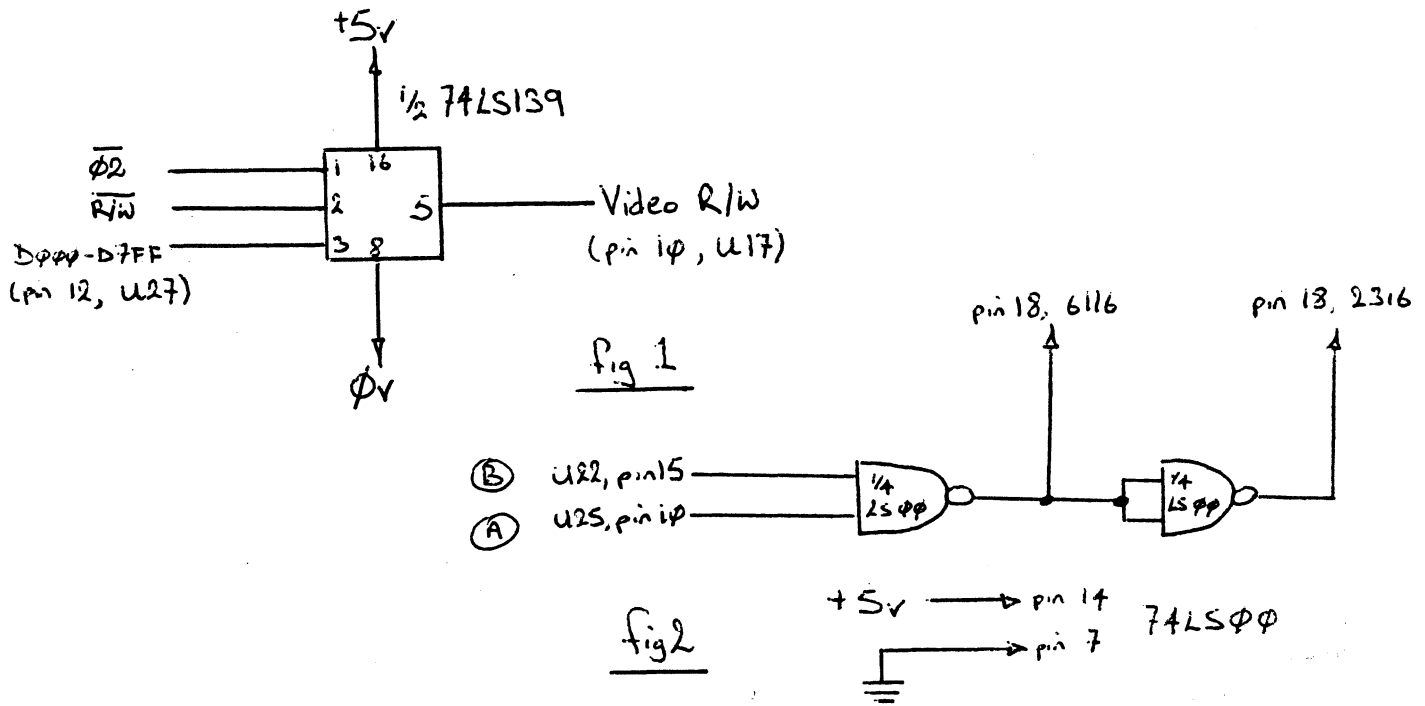
## THE TASAN VIDEO BOARD - SOME "FIXES"

This is probably going to be the last article that I shall write on the TASAN Video board, so I will devote it to some alterations (both to it, and to my programmable character generator - see Vol.2 No.9 newsletter).

Many of you who have built the Video board have complained about random characters occurring here and there on the screen. The first method I would suggest for fixing this problem is to replace IC 26 (a 74LS03) with a 74LS38. If this doesn't seem to fix the problem then you should piggy-back 74LS139 on to the 74LS139 on the board (IC 48) soldering pins 1, 2, 8 and 16; bend all other pins out a little so they don't make contact with the pins beneath them. Solder a length of wire wrap wire from pin 3 of your new chip to pin 12 or 13 of IC 27 (a 74LS20). Solder a length of wire wrap wire from pin 5 of your new chip to pin 10 of IC 17 (a video 2114); cut the track from pin 12 of IC 16 at the pin. This should fix the problem.

The cause of this problem is that when the computer is accessing the screen RAM, it is given priority over all the video signals, including the P/W line to the video RAMs, so, if you write data to the screen, the R/W line goes low to indicate writing; but when the computer has finished writing to the screen, priority over the video signals is given back to the scanning circuitry. Unfortunately, the 74LS157s which do this, appear to be a little slow and the R/W line may still be in the write state, which causes a character to appear at a random address on the screen.

The new video R/W circuitry looks like fig.1



Some Video boards have problems where the characters in the 32 X 32 format "break up" or are otherwise "messy". This is easily cured by replacing U5 - a 74LS157 with a standard (not LS) 74157.

If your screen is plagued by black lines running through it - especially when colour is enabled. This is cured by replacing the 8T28s with faster ones, (just replace them, starting from the right hand end until the problem is fixed).

There are two problems with my Programmable character generator, firstly on my circuit diagram, I have labeled pin 20 of the 6116 and 2316 wrongly. Pin 20 of the 6116 should be grounded, and pin 20 of the 2316 should have the strap option, which I have shown coming from pin 20 of the 6116.

Secondly, the circuitry to pin 18 of the 6116 and 2316 is slightly wrong. The new circuitry should look like fig.2

Paul Dodd

## TASNET

The Tasmanian Education Department Network (TASNET) through the Elizabeth Computer Centre (E.C.C.) services Schools, Colleges and Administrative Departments throughout Tasmania with a number of computers in various regions linked together by synchronous telephone line (4800 baud) or local asynchronous lines (1 megabaud in Hobart).

TASNET currently consists of one Digital VAX 11/780, three PDP 11/70s and two 11/34s at the following locations:

HOBART	:	1 only VAX 11/780	.....	2 Mbyte
		2 only PDP 11/70	.....	$\frac{1}{2}$ Mbyte each
LAUNCESTON	:	1 only PDP 11/70	.....	$\frac{1}{2}$ Mbyte
DEVONPORT	:	1 only PDP 11/34	.....	$\frac{1}{2}$ Mbyte
BURNIE	:	1 only PDP 11/34	.....	$\frac{1}{2}$ Mbyte

There is approximately 2000 Mbytes of on-line storage available with off-line magnetic tape backup available at each computer - the whole system is backed up daily.

All High Schools and Colleges now have on-line facilities - usually by either 300 baud dedicated lines or 300 baud dial-up modems connected to V.D.U.s and/or printers (LA 34 or LA 36 usually) with on-line facilities for high speed line printers at each computer node. Many schools also have APPLE microcomputers with 48K memory and one disk drive. Software has been developed to allow down loading of Applesoft programs stored on the system and APPLES to be used as TASNET terminals.

Uses to which the system is put include:

1. Administration
2. Computer Studies courses at Years 9, 10, 11, 12
3. Computer Assisted Learning
4. Computer Awareness Courses at Years 7, 8.
5. Library Circulation System

In the Administration applications area, teacher support is offered through:

1. TASCIS an on-line cataloguing data base for school and college libraries - it now has over 100,000 books catalogued and library cards for these may be requested for printout on a special terminal.
2. THE TASMANIAN MEDIA CENTRE has over 10,000 items available for loan (films, tapes, videos, kits) and information on these items is available on-line to all schools (includes synopsis, suitability, availability) and schools can check their bookings (own only).
3. ABC REPORTER All schools TV and RADIO programs are catalogued on-line and times, suitability and synopsis for each is available as well as coming highlights and alterations.

Many schools have computerised their administration with various listings of classes, addresses, subjects and groupings, census information etc. readily available.

The system has available BASIC, EXTENDED BASIC, PASCAL, COBOL, FORTRAN, LOGO, INTERFACING TO TURTLES AND A FILE MANAGEMENT PACKAGE. Computer assisted learning, games and other general purpose programs are available including assessment, survey and census data information.

Library circulation software has been developed in TASMANIA and is used for on-line circulation control in some branches of the State Library and in some schools and colleges.

APPLE microcomputers are supported and many programs have been written by E.C.C. staff and by teachers and are available for sale to interested groups and persons in other states. Very recently B.B.C. Protons have been recommended as the new supported computer for the next five years, although APPLES will continue to be supported. If enough interest is shown further information on APPLES could be made available in a later issue.

Graeme D. Reardon  
Senior Master Maths/Science  
Cressy District High School  
Tasmania

## THE DEBUGGER

The debugger is an extremely useful tool for persuading machine code routines to work. It allows the user to single step his program, returning control to a trace routine each time.

The key to the debugger lies in the hardware. In fact, all the debugger consists of is a 3 input NAND gate. (Makes you wonder why all micros don't have it.) The 6502 chip outputs a pulse on the SYNC pin each time it begins an instruction. If this is fed back to the NMI input, an interrupt would be generated each time an instruction is executed. Clearly, this would be useless as the NMI routine would also be traced and the system would wind itself up in knots. However, if this signal was passed through a NAND gate, (the SYNC output sends a positive going pulse, NMI is active low) the chip select line from the trace routine ROM could disable the debugger while executing the trace program. The third input of the NAND gate is simply connected to a flip-flop which is toggled by the 'DBG ON' and 'OFF' switches on the SYM keypad.

The trace routine itself can be as simple or as complex as you like. SUPERMON merely displays the contents of the program counter and accumulator, delays (or stops if TV at \$A656=0) and resumes execution of the program. There are, however, a few problems involved with user trace routines.

A user trace routine will be located in RAM or ROM outside the SUPERMON ROM. Therefore, the chip select signal which disables the debugger is ineffective. The SYM solves this problem by allowing the debug switches to be controlled from software. By inserting links W-24 and X-25, VIA#3 at \$AC00 can now control the debugger. These links are located just above the speaker.

The initialization routine of your new debugger must put \$80C0 in the NMI vector SA67A and the address of the actual trace routine in \$A674. When the NMI is raised, SUPERMON will call ACCESS, save the registers, switch off the debugger and jump to your tracer which should then print the registers, memory or whatever, and finish with a -

4C CD80                      JMP TRACON

to switch on the debugger and resume execution. Some condition should be included to halt execution or the debugger will trace forever or until the program returns to SUPERMON.

There are a few interesting points to note with regards to the debugger.

1. DO NOT USE THE 'BRK' INSTRUCTION. This instruction in a user program will interfere with the debugger's operation.
2. The debugger can trace programs in ROM. Many similar systems insert BRK instructions in the user program for single stepping which requires the program to be in RAM. I have successfully single stepped BASIC with the SYM debugger.
3. The chip select disabling capability means that the debugger will not bother single stepping SUPERMON subroutine calls.
4. The NMI input to the 6502 can not be used for any other purpose without hardware modifications.

Try a simple trace routine at first to get the feel for the debuggers operation before attempting a major trace program.

NEXT MONTH - THE SPEAKER

In the final SYM-POSIUM article, we will look at the small piezo-electric speaker on the SYM and how it can be used to play music or generate sound effects. We will also look at other methods of sound generation.

Brian Campbell

## THE MEETING WAS KAOS

RABBLE EXPANSION BOARD: Bill Chilcott had his new board on show at the meeting. Details on availability and price are listed elsewhere in this newsletter.

65C02: Rockwell has released their new 65C02 5 MHz CMOS microprocessor. The new chip will contain an extra 27 instructions over the Mos-Tech 6502. Rockwell is under licence to Mos-Tech.

SOFT FRONT PANEL: A part of the GTBUG package written by Tony Durrant is his new soft front panel. This new feature will single step a program, change memory, disassemble instructions into mnemonics and breakpoint can also be set. Tony will not release GTBUG for sale until he is completely satisfied with it, so I'm afraid we will all have to patient for a while longer.

FORTH: mention of a new local newsletter based around Forth is in the pipeline and should interest most Forth users.

TAPE LIBRARY: John Whitehead, known for his numerous articles in the KAOS newsletters, has volunteered to run the Cassette Library from now on. John's phone number is 763 5983.

Those were the main points of the meeting though much more information is exchanged during the ragchew before and after the meeting.

73's from Melbourne

Rod Drysdale VK3BYU

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### SPECIAL DEALS FOR KAOS MEMBERS FROM COMP- SOFT

Rabble Expansion Board	Board only	\$85.00
	Full Kit	\$305.00
	Assem. & Tested	\$380.00

COMING SOON: BMC High Resolution Green Screen, 15MHz band width \$195.00

Microline 80 Printer	still available	at	\$535.00
2732 EPROM	good price		\$5.50

Special price for 1 month till the 30th September.

TASAN VIDEO BOARD	64X32 screen format	Board only	\$40.00
		Partial Kit	\$115.00
		Full Kit	\$130.00

### RABBLE EXPANSION BOARD

The Rabble Ozi Expansion Board is designed to suit Superboard 11 and C1P computer cases. It provides for 32K of 6116LP3 RAM, addressed at locations \$0000 - \$7FFF. Facilities are also provided for a further 32K of ROM at locations \$8000 - \$FFFF, using 2732s. The Board is strappable to suit 2532s, 2716s and also 2764s.

The Rabble Ozi Expansion Board has an OHIO compatible Floppy Disk Controller for 5 $\frac{1}{4}$ " or 8" drives, complete with data separator and motor controller. Other features are a Real Time Clock, a PIA and VIA to provide I/O connections for a parallel printer. Last is the dual Programmable Sound Generators providing 6 channels of sound which can provide a simple noise such as a gunshot or explosion.

The current consumption of this Board is less than 1 Amp at 5V (current varies with loads on I/Os).

Available from:

Rabble Ozi Computers  
P.O. Box 781  
Shepparton Vic. 3630

COMP-SOFT  
235 Swan St  
Richmond Vic. 3121

Data Parts  
5 Naomi St  
Shepparton Vic. 3630

65U DOS RAMBLINGS  
*by David Tasker*

A few months ago I obtained a copy of 65U for my much modified C1. After conferring with a couple of the more active disk users in KAOS who suggested my time might be better used designing more hardware, I decided to try and unravel some of the differences between the 65D and 65U operating systems.

65U is best understood as a core or kernal based system. i.e. There is a central operating system which is accessed by the high level languages as required. We can send commands to the core from the Extended Monitor, Assembler or Basic. Typically, if we were in Basic and we wanted to load a program we would type: DISK!"LOAD FILNAM" We could then list and modify the program and RUN or SAVE it. If we wanted to SAVE it under a different filename, we would have to CREATE a place for it on the disk with a program called CREATE. 65U would SAVE the program by Track and sector.

65U is primarily a BASIC system. Commands to the DOS for load could be, LOAD"FILNAM" if the original program was read/write access (R/W) then no password is needed. However, we can with 65U specify at the time we CREATED the program location, assign a 4 character password. e.g. LOAD"FILNAM","PASS",10 where PASS is our password and 10 is the line number to enter the program.

With 65D you might consider that all the DOS commands are a subset of the BASIC word DISK!" where as in 65U the DOS commands are a part of BASIC. Another example. Say we currently booted our system in from Drive A. (This would normally be the case.) Part way through our program we wanted to get a DATA file from the Disk loaded in Drive B. The syntax for 65D and 65U including line numbers is shown to compare.

65D	65U
10 DISK!"SE B"	10 DEV"B"
20 DISK OPEN,6, "FILNAM"	20 OPEN "FILNAM",1

The 6 in the 65U is a buffer in memory called device 6. Anything in this BUFFER can be inputed into the program (INPUT 6) and we can print to this buffer (PRINT 6). There are two buffers that we can use (6 and 7).

With 65U the 1 indicated a channel (channel 1). We can input (INPUT%1) and print (PRINT%1) also CLOSE device or channel number can be used. This will cause any new or changed data to be written back to the disk file. In 65U you can have up to 8 channels opened at any one time.

65U does not have a kernal or core that you access. There is no Extended Monitor or Assembler. If you want to access Assembler or 65D files, then 65U does have a program called LOAD32 (for 32K systems) and LOAD48 (for 48K systems) which gives you a kernal like operating system for accessing 65D files.

65U also requires more operating memory. The DOS is 24K long and needs 8K of work-space. So a minimal system is 32K. Quite a few very powerful BASIC keywords for file handling.

*Continued next month*

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FOR SALE

AIM 65 with 12K of Assembler Monitor in ROM, 36K memory, Digital Cassette Recorder, EPROM programmer, 60 I/O lines, Printer on board. All one integrated unit, in one case with power supply. Used for commercial development. \$590.00 or ono

Contact Grahame Younghusband

SUPERBOARD II, 8K RAM, DABUG III, 300 and 600 baud, 1-2 MHz, Custom made case, Modulator, Power Supply, Manuals - standard and additional (several) Approximately \$200 in software, includes Editor/Assembler, Extended Monitor, utilities and many games. REDUCED FOR A QUICK SALE \$400.00 or near offer.

Contact Alex Koehler